# Evaluating the $\mathrm{M}|\mathrm{D}| \infty$ Queue Busy Cycle Distribution 

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# Evaluating the $M|D| \infty$ Queue Busy Cycle Distribution 

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#### Abstract

Given the busy period major importance in queuing systems, it is also relevant to study the busy cycle. In this work some interesting results on $M|G| \infty$ queue system busy cycle distribution are presented. They are emphasized for the $\mathrm{M}|\mathrm{D}| \infty$ queue system and a numerical method to compute the $\mathrm{M}|\mathrm{D}| \propto$ queue system busy cycle distribution function is presented.


Keywords: $\mathrm{M}|\mathrm{D}| \propto, \mathrm{M}|\mathrm{G}| \infty$, busy cycle, distribution function.

## 1 Introduction

A queue system busy period is a period that begins when a customer arrives at the system finding it empty, ends when a customer abandons the system letting it empty and, throughout its progress, there is always at least one customer present. An idle period followed by a busy period is a busy cycle.

In the $\mathrm{M}|\mathrm{G}| \infty$ queue system the customers arrive according to a Poisson process at rate $\lambda$, receive a service which time length is a positive random variable with distribution function $G($.$) and mean \alpha$ and, when they arrive, each one finds immediately an available server. Each customer service is independent from the other customers' services and from the arrivals process. The traffic intensity is $\rho=\lambda \alpha$.

Call $I, B$ and $Z$ the time length random variable of the idle period, the busy period and the busy cycle respectively; $i(t), b(t)$ and $z(t)$ are the correspondent probability density functions and $I(t), B(t)$ and $Z(t)$ the distribution functions.

[^0]
## 2 General Results

Evidently $Z=I+B$ and being I and B independent, see [2], the distribution of $Z$ is the $I$ and $B$ distributions convolution. Then, being $\bar{Z}(s), \bar{I}(s)$ and $\bar{B}(s)$ the $\mathrm{Z}, B$ and $I$, respectively, Laplace transforms

$$
\begin{equation*}
\bar{Z}(s)=\bar{I}(s) \bar{B}(s) \tag{2.1}
\end{equation*}
$$

where

$$
\begin{equation*}
\bar{I}(s)=\frac{\lambda}{\lambda+s} \tag{2.2}
\end{equation*}
$$

as it happens for any queue with Poisson arrivals process and

$$
\begin{equation*}
\bar{B}(s)=1+\lambda^{-1}(s- \tag{2.3}
\end{equation*}
$$

$\left.\frac{1}{\int_{0}^{\infty} e^{-s t-\lambda} \int_{0}^{t}[1-G(v)] d v} d t\right)$
see again [2].
Consequently

$$
\begin{equation*}
E\left[Z^{n}\right]=\sum_{p=0}^{\infty}\binom{n}{p} \frac{p!}{\lambda^{p}} E\left[B^{n-p}\right], n=1,2, \ldots \tag{2.4}
\end{equation*}
$$

where, see [6],

$$
\begin{gather*}
E\left[B^{n}\right]=(-1)^{n+1}\left\{\frac{e^{\rho}}{\lambda} n C^{(n-1)}(0)-e^{\rho} \sum_{p=1}^{n-1}(-1)^{n-p}\binom{n}{p} E\left[B^{n-p}\right] C^{(p)}(0)\right\} \\
n=1,2, \ldots \tag{2.5}
\end{gather*}
$$

and

$$
\begin{equation*}
C^{(n)}(0)=\int_{0}^{\infty}(-t)^{n} e^{-\lambda \int_{0}^{t}[1-G(v)] d v} \lambda(1-G(t)) d t, n=01,2, \ldots \tag{2.6}
\end{equation*}
$$

So

$$
\begin{equation*}
E[Z]=\frac{e^{\rho}}{\lambda} \tag{2.7}
\end{equation*}
$$

does not depend on the service time distribution form, except for its mean $^{2}$. But $E\left[Z^{n}\right], n \geq 2$ depend on the whole service time distribution structure.

For the $M|D| \propto$ queue system - constant $^{3}$ service times with value $\alpha$ -

$$
\begin{equation*}
\bar{B}(s)=1+\lambda^{-1}\left(s-\frac{(s+\lambda) s}{\lambda e^{-(s+\lambda) \alpha}+s}\right) \tag{2.8}
\end{equation*}
$$

obtaining, by Laplace transform inversion, see $[3]^{4}$,

$$
\begin{equation*}
b(t)=\sum_{n=0}^{\infty}\left(\frac{d}{d t} \frac{c(t)}{e^{-\rho}}\right) *\left(\frac{d}{d t} \frac{1-d(t)}{1-e^{-\rho}}\right)^{* n} e^{-\rho}\left(1-e^{-\rho}\right)^{n} \tag{2.9}
\end{equation*}
$$

where $\frac{c(t)}{e^{-\rho}}=\left\{\begin{array}{l}0, t<\alpha \\ 1, t \geq \alpha\end{array}=G(t)\right.$ and $\frac{1-d(t)}{1-e^{-\rho}}=\left\{\begin{array}{c}\frac{1-e^{-\lambda t}}{1-e^{-\rho}}, t<\alpha \\ 1, t \geq \alpha\end{array}\right.$.
Then

$$
\begin{equation*}
\bar{Z}(s)=1-\frac{s}{\lambda e^{-(s+\lambda) \alpha}+s} \tag{2.10}
\end{equation*}
$$

and

$$
\begin{equation*}
z(t)=\left(\lambda e^{-\lambda t}\right) * b(t), t \geq 0 \tag{2.11}
\end{equation*}
$$

Still

$$
\begin{gather*}
C^{(0)}(0)=1-e^{-\rho} \\
C^{(n)}(0)=-e^{-\rho}(-\alpha)^{n}-\frac{n}{\lambda} C^{(n-1)}(0), n=1,2, \ldots \tag{2.12}
\end{gather*}
$$

see again [6].

[^1]
## 3 The $\boldsymbol{M}|\boldsymbol{D}| \infty$ Queue Busy Cycle Distribution Function

The expression (2.11) for $z(t)$, in the former section, allows the busy cycle distribution structure interpretation for the $M|D| \infty$ queue. But it fails in the task of presenting an easy expression for the distribution function $Z(t)$ computation.

This may be done, for example, with an algorithm created by Platzman, Ammons and Bartholdi III, see [1] ${ }^{5}$, that allows the distribution functions computation since the correspondent Laplace transform in round form is known, as it is now the case, remember (2.10). Unhappily the same does not happen for other $M|G| \infty$ systems what inhibits the use of this algorithm.

[^2]\[

$$
\begin{equation*}
P[X \geq t+\Delta t]-\Delta p \leq \tau \leq P[X>t-\Delta t]+\Delta p \tag{3.1}
\end{equation*}
$$

\]

Platzman, Ammons and Bartholdi III suggest doing

$$
\begin{equation*}
\tau=\frac{U-t+\Delta t}{U-L+2 \Delta t}+\sum_{n=1}^{N} \frac{\alpha^{n^{2}}}{\pi n} \operatorname{im}\left\{\left(\beta^{n}-\gamma^{n}\right) L(j \omega n)\right\} \tag{3.2}
\end{equation*}
$$

where $K=\log \frac{2}{\Delta p}, \quad \mathrm{D}=\frac{\Delta t}{\sqrt{2 K}}, \omega=\frac{2 \pi}{U-L+2 \Delta t}, \mathrm{~N}=\left[\frac{2 K}{\omega \Delta t}\right]$, being $[\cdot]$ the characteristic of a real number, $\quad \alpha=e^{-D^{2} \frac{\omega^{2}}{2}}, \quad \beta=e^{j(U+\Delta t) \omega}, \quad \gamma=e^{j t \omega}, U$ and $L$ are numbers such that $1-$ $P[L \leq X \leq U] \ll \Delta p, j=\sqrt{-1}$ and $\operatorname{im}(\cdot)$ designates the imaginary part of a complex number. $L(j \omega n)$ is the Laplace transform value in $j \omega n$. They demonstrate that the approximation so defined fulfills the condition (3.1).

The algorithm implementation, for details see [4], is computationally performed through a FORTRAN program, see [5], and the results of some experiences are presented in the Annex.

The values of $\alpha, \lambda, \Delta t$ and $\Delta p$ must be specified and the values of $t$ for which the values of $Z(t)$, called $Z^{c}(t)$, are wanted.

As for the goodness of the obtained results, it is tested computing the errors of $E\left[Z^{c}\right]$ and $V A R\left[Z^{c}\right]$, computed after them, in relation with the true values of $E[Z]$ and $\operatorname{VAR}[Z]$ that are available for this queue system. The exception is the first experience where, with $\alpha=0$, the situation is the one of a pure Poisson process. So, the results obtained ( $2^{\text {nd }}$ column in Table 1) are compared with the Poisson process ones ( $3^{\text {rd }}$ column in Table 1). Generally, the $Z^{c}$ values fit well.

## References

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## ANNEX

Table 1. Experience 1: $\alpha=0, \lambda=1, \Delta \mathrm{t}=0.01$ and $\Delta \mathrm{p}=0,001$

| t | $Z^{c}(\mathrm{t})$ | Poisson Process |
| :--- | :--- | :--- |
| 0 | 0.00020928263 | $0.000 \ldots$ |
| 0.5 | 0.39354845 | 0.39346934 |
| 1 | 0.63201874 | 0.632120559 |
| 1.5 | 0.77676630 | 0.77686984 |
| 2 | 0.86456292 | 0.864664717 |
| 2.5 | 0.91781115 | 0.917915001 |
| 3 | 0.95011103 | 0.95021212932 |
| 3.5 | 0.96969878 | 0.969802617 |

Table 2. Experience 2: $\alpha=1, \lambda=1, \Delta t=0.01$ and $\Delta \mathrm{p}=0,001$

| t | $Z^{c}(\mathrm{t})$ | t | $Z^{c}(\mathrm{t})$ |
| :---: | :---: | :---: | :---: |
| 0.5 | 0.00070788896 | 4.5 | 0.89332950 |
| 1 | 0.00078194999 | 5 | 0.92884773 |
| 1.5 | 0.18467983 | 5.5 | 0.95303684 |
| 2 | 0.36851909 | 6 | 0.96932029 |
| 2.5 | 0.53561949 | 6.5 | 0.98016983 |
| 3 | 0.66881525 | 7 | 0.98734205 |
| 3.5 | 0.76919734 | 7.5 | 0.99205017 |
| 4 | 0.84198290 | $\operatorname{VAR}[Z]=1.9444392442$ |  |
| $\begin{gathered} E[Z]=2.718281829 \\ E\left[Z^{c}\right]=2.605018789 \\ \varepsilon=4 \% \end{gathered}$ |  |  |  |
|  |  | $\begin{aligned} \operatorname{VAR}\left[Z^{c}\right] & =1.875647136 \\ \varepsilon & =3.5 \% \end{aligned}$ |  |

Table 3. Experience 3: $\alpha=1, \lambda=2, \Delta t=0.01$ and $\Delta \mathrm{p}=0,001$

| t | $Z^{c}(\mathrm{t})$ | t | $Z^{c}(\mathrm{t})$ |
| :--- | :--- | :--- | :--- |
|  |  |  |  |
| 0.5 | 0.00038790601 | 7.5 | 0.92047894 |
| 1 | 0.00045109048 | 8 | 0.93518191 |
| 1.5 | 0.13572108 | 8.5 | 0.94718128 |
| 2 | 0.27099844 | 9 | 0.95697385 |
| 2.5 | 0.39718168 | 9.5 | 0.96496373 |
| 3 | 0.50513958 | 10 | 0.97148519 |
| 3.5 | 0.59509700 | 10.5 | 0.97680729 |
| 4 | 0.66922503 | 11 | 0.98115152 |
| 4.5 | 0.72997826 | 11.5 | 0.96469930 |
| 5 | 0.77964925 | 12 | 0.98759257 |
| 5.5 | 0.82022225 | 12.5 | 0.98995178 |
| 6 | 0.85335999 | 13 | 0.99188309 |
| 6.5 | 0.88039940 | 13.5 | 0.99344980 |
| 7 | 0.92047130 | 14 | 0.99473917 |


| $E[Z]=3.69452805$ | $\operatorname{VAR}[Z]=6.260481408$ |
| :---: | :---: |
| $E\left[Z^{c}\right]=3.606224458$ | $\operatorname{VAR}\left[Z^{c}\right]=5.358674148$ |
| $\varepsilon=2.4 \%$ | $\varepsilon=14 \%$ |

Table 4. Experience 1: $\alpha=2, \lambda=1, \Delta t=0.01$ and $\Delta p=0,01$

| t | $Z^{c}(\mathrm{t})$ | t | $Z^{c}(\mathrm{t})$ |
| :--- | :--- | :--- | :--- |
| 0.5 | 0.00039526703 | 14 | 0.90255320 |
| 1 | 0.00039531649 | 14.5 | 0.91201680 |
| 1.5 | 0.00039744257 | 15 | 0.92056465 |
| 2 | 0.00042999497 | 15.5 | 0.92828899 |
| 2.5 | 0.0068082088 | 16 | 0.93526571 |
| 3 | 0.13566480 | 16.5 | 0.94157290 |
| 3.5 | 0.20333376 | 17 | 0.94726365 |
| 4 | 0.27105104 | 17.5 | 0.95241045 |
| 4.5 | 0.33643096 | 18 | 0.95705801 |
| 5 | 0.39722785 | 18.5 | 0.96125179 |
| 5.5 | 0.45344632 | 19 | 0.96504825 |
| 6 | 0.50523263 | 19.5 | 0.96847575 |
| 6.5 | 0.55233818 | 20 | 0.97157025 |
| 7 | 0.59518069 | 20.5 | 0.97437018 |
| 7.5 | 0.63407224 | 21 | 0.97689431 |
| 8 | 0.66930794 | 21.5 | 0.97917509 |
| 8.5 | 0.70120662 | 22 | 0.98124003 |
| 9 | 0.73005634 | 22.5 | 0.98309797 |
| 9.5 | 0.75615197 | 23 | 0.98477888 |
| 10 | 0.77973318 | 23.5 | 0.98630297 |
| 10.5 | 0.80105113 | 24 | 0.98767584 |
| 11 | 0.82031202 | 24.5 | 0.98891764 |
| 11.5 | 0.83771467 | 25 | 0.99003869 |
| 12 | 0.85343867 | 25.5 | 0.99104917 |
| 12.5 | 0.86764937 | 26 | 0.99196279 |
| 13 | 0.88047999 | 26.5 | 0.99279278 |
| 13.5 | 0.89207541 | 27 | 0.99353820 |
| $E[Z]=7.389056099$ |  | $V A R[Z]=25.04192563$ |  |
| $E\left[Z^{c}\right]=7.200722486$ | $V A R\left[Z^{c}\right]=20.69584719$ |  |  |
| $\varepsilon=2.5 \%$ |  |  |  |
| $\varepsilon$ | $\varepsilon=17 \%$ |  |  |


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[^1]:    ${ }^{2}$ In these circumstances it is usual to say that it is insensible to the service time distribution.
    ${ }^{3}$ That is: Deterministic service times.
    ${ }^{4} *$ is the convolution operator.

[^2]:    5 It is generally said that an algorithm is "accurate" if it looks for solving a problem "close" to the one that is supposed to solve. An algorithm is "precise" if it gets a solution "close" to the one of the problem that it is trying to solve. .More concretely, being $\Delta t(\Delta t>0)$ the accuracy and $\Delta p\left(0<\Delta p<\frac{1}{2}\right)$ the precision required, the approximation $\tau$ of $P[X>t]$ must satisfy the condition

